

# CS 263: Counting and Sampling

Nima Anari



slides for

## Continuous Time

# Review

## $\phi$ -entropy

For function  $\phi$  and  $f : \Omega \rightarrow \mathbb{R}$  define

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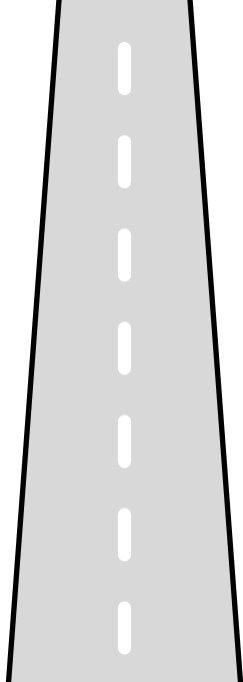
- ▶ Eigvals are  $\mathbb{E}_{\pi}[\chi]$

# Fourier Analysis

- ▶ Characters
- ▶ Examples
- ▶ Relaxation time

# Continuous Time

- ▶ Functional analysis in continuous time
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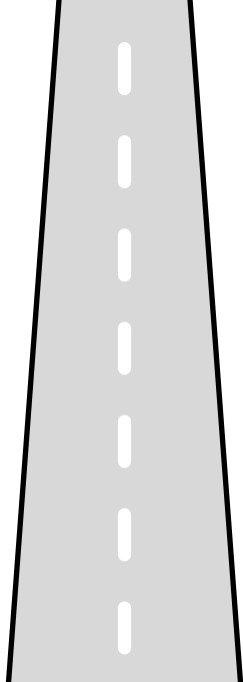


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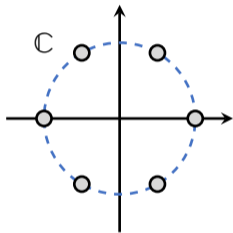


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$$\chi(x) = \omega^x$$

for  $\omega$  an  $n$ -th root of unity.

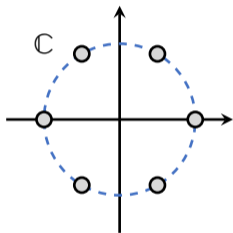


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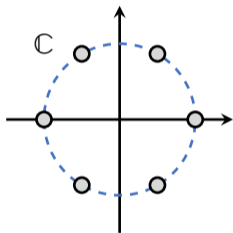


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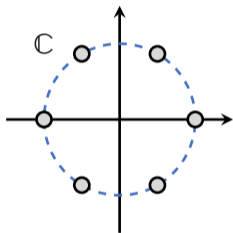
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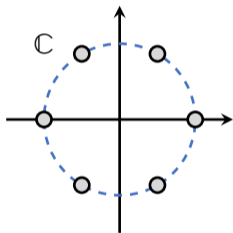
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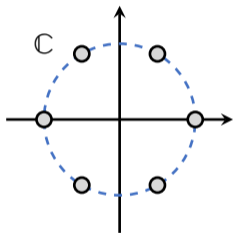
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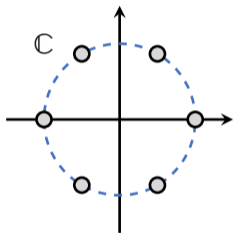
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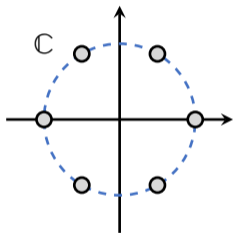
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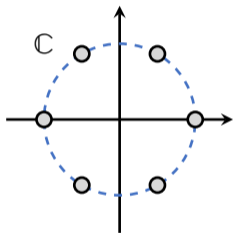
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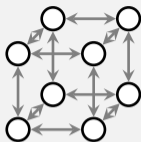
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- ▶ **Mixing**: largest  $|\cdot|$  of an eig?

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▶  $1_i$  w.p.  $1/2n$



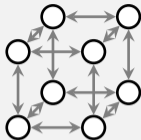


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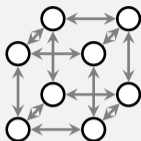
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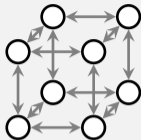
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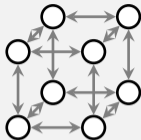
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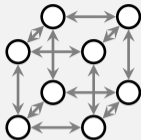
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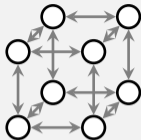
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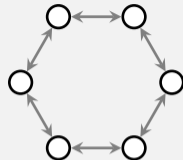
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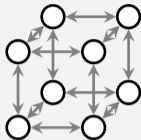
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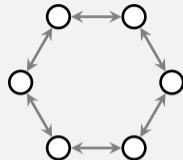
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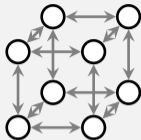
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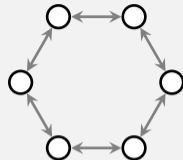
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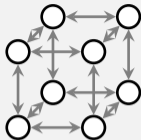
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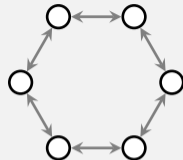
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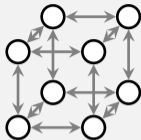
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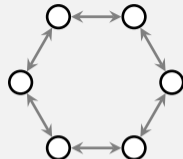
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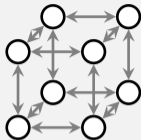
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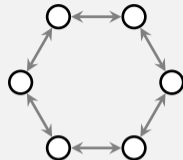
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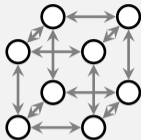
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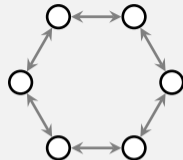
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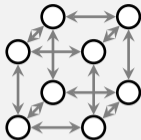
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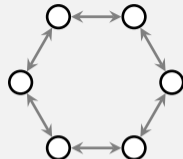
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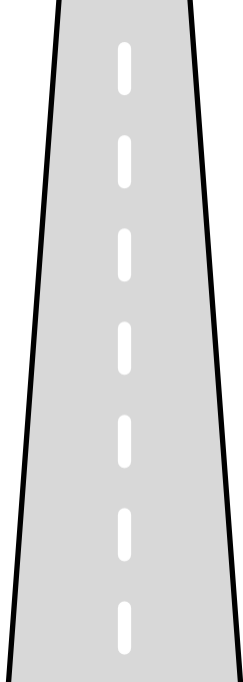
▶ Note: going back from  $\lambda_2$  to  $t_{\text{mix}}$  gives us non-tight bound of  $O(n^2)$ . 😞

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- ▶ Characters
- ▶ Examples
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- ▶ Functional analysis in continuous time
- ▶ Dirichlet form

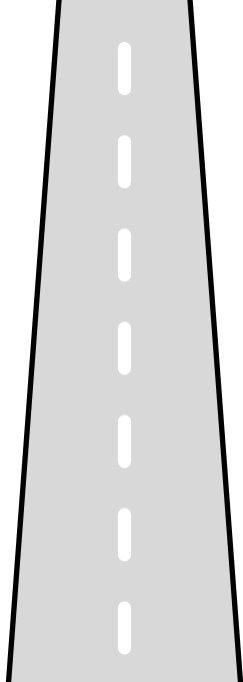


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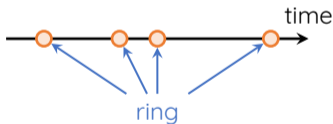
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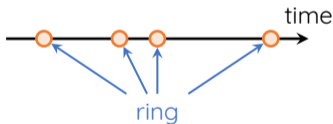
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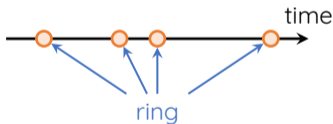
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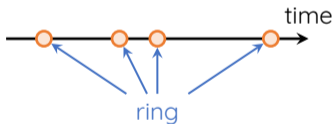
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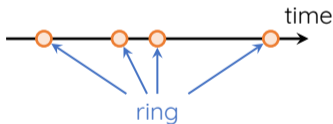
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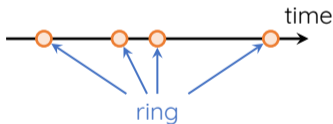
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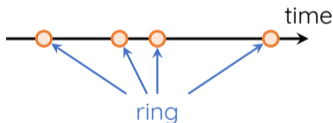
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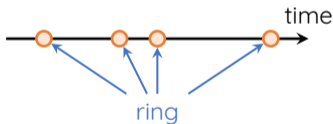
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- ▶ Easier because of **Dirichlet form!**

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Define  $\mathcal{E}(f, g)$  for functions  $f, g : \Omega \rightarrow \mathbb{R}$  as

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Poincaré:  $2 \mathcal{E}(f, f) \geq \rho \text{Var}[f]$

MLSI:  $\mathcal{E}(f, \log f) \geq \rho \text{Ent}[f]$

Just need to lower bound  $\varepsilon$