

## Review

- Path Coupling:  $W(1_x P, 1_{x'} P) \leq (1-c) W(1_x, 1_{x'})$

$\downarrow$   
only adjacent states  
\* Coloring:  $q > 2\Delta \Rightarrow t_{\text{mix}}^{\text{(Metropolis)}} = O(n \lg n)$

- Dobrushin:  $\mu$  on  $\Omega_1 \times \dots \times \Omega_n$

$$I[j \rightarrow i] = \max_{\{x_i, x'_i\}} \left\{ \frac{d(\mu_i | x_i, \mu_i | x'_i)}{TV} \mid x \sim_j x' \right\}$$

Thm: If columns of  $I$  sum to  $\leq 1-\delta$   
 $\Rightarrow t_{\text{mix}}(\text{Glauber}) = O(n \lg n / \delta)$

\* Coloring:  $q > 2\Delta \Rightarrow t_{\text{mix}} = O(n \lg n)$

\* Hardcore:  $\lambda \leq \frac{1-\delta}{\Delta} \Rightarrow t_{\text{mix}} = O_{\delta}(n \lg n)$

\* Ising:  $\sum_i |\beta_{ij}| \leq 1-\delta \Rightarrow t_{\text{mix}} = O_{\delta}(n \lg n)$

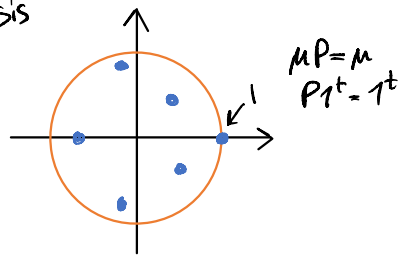
- Dobrushin  $\dagger$ :  $\lambda_{\max}(I)$

$$C \in \mathbb{R}_{>0}^n \text{ with } CI \leq (1-\delta)C$$

$$\Rightarrow W(\nu P, \nu' P) \leq (1-\frac{\delta}{n}) W(\nu, \nu')$$

$\rightarrow$   $c$ -weighted Hamming distance

## - Spectral Analysis



- Time-Reversible: eigenvalues are real

$$-1 \leq \lambda_n \leq \dots \leq \lambda_2 \leq \lambda_1 = 1$$

- Spectral Gap:  $1-\lambda_2$  or  $1-\max(|\lambda_2|, |\lambda_n|)$

- Relaxation Time:  $t_{\text{rel}} := \frac{1}{1-\lambda_2}$  ← a proxy for  $t_{\text{mix}}$

## Plan for Today

- Relationship between  $t_{\text{mix}}, t_{\text{rel}}$

- Intro to functional analysis

- Fourier analysis

# Eigenvalues & Mixing Time

Suppose  $vP = \lambda v$  for  $\lambda \neq 1 \in \mathbb{C}$ ,  $v \in \mathbb{C}^n$ .

$$\lambda v 1^t = v P^t = v 1^t \Rightarrow v 1^t = 0$$

$$\text{Def: } \|v\|_1 = \sum_i (|\operatorname{Re}(v_i)| + |\operatorname{Im}(v_i)|)$$

Claim:  $t \geq t_{\text{mix}}(\epsilon) \Rightarrow \|v P^t\|_1 \leq O(\epsilon) \|v\|_1$

Proof: Let  $v := \alpha v_1 - \alpha v_2 + i\beta v_3 - i\beta v_4$  for

$v_1, v_2, v_3, v_4$  dists:  $v_i \in \mathbb{R}_{\geq 0}^n$ ,  $v_i 1^t = 1$

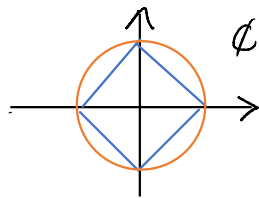
$v_1, v_2$  disjoint support  
 $v_3, v_4$  disjoint support  $\Rightarrow \|v\|_1 = 2\alpha + 2\beta$

$$v P^t = \alpha (v_1 \overset{M}{P^t} - v_2 \overset{M}{P^t}) + i\beta (v_3 \overset{M}{P^t} - v_4 \overset{M}{P^t})$$

$$\Rightarrow \|v P^t\|_1 \leq \alpha \cdot O(\epsilon) + \beta \cdot O(\epsilon) = O(\epsilon) \cdot \|v\|_1$$

$$\text{Def: } \|v\|_1' = \sum_i |v_i| \downarrow \sqrt{\operatorname{Re}(v_i)^2 + \operatorname{Im}(v_i)^2}$$

$$\|v\|_1' \stackrel{\sqrt{2}}{\geq} \|v\|_1 \quad \text{but} \quad \|\lambda v\|_1' = |\lambda| \|v\|_1'$$



Now if  $t \geq t_{\text{mix}}(\epsilon)$  we have

$$|\lambda|^t \|v\|_1' = \|v P^t\|_1' \leq O(\epsilon) \cdot \|v\|_1' \\ \Rightarrow |\lambda|^t = O(\epsilon)$$

\* Corollary:  $|\lambda| < 1 \Leftarrow$  Holds even when for eigenvectors other than  $\mu$  if ergodic.

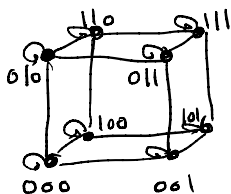
$$* \text{Corollary: } 1 - |\lambda| \geq \Omega\left(\frac{1/\epsilon + \text{const}}{t_{\text{mix}}(\epsilon)}\right)$$

$$|\lambda| \leq 1 - \Omega\left(\frac{1}{t_{\text{mix}}}\right) \leftarrow \text{usually not tight}$$

Example. (Hypercube)

$$t_{\text{mix}}(\varepsilon) = O(n \lg n + n \lg \frac{1}{\varepsilon})$$

$$\Rightarrow t_{\text{rel}} = O\left(\frac{n \lg n + n \lg \frac{1}{\varepsilon}}{\text{const} + \lg \frac{1}{\varepsilon}}\right) \xrightarrow{\varepsilon \rightarrow 0} t_{\text{rel}} = O(n)$$



Example. (Dobrushin)

Suppose  $\sum_i |I_i[j \rightarrow i]| \leq 1 - \delta$

Glauber

$$\Rightarrow t_{\text{mix}} = O\left(\frac{n \lg n}{\delta} + \frac{n \lg \frac{1}{\varepsilon}}{\delta}\right)$$

$$\Rightarrow t_{\text{rel}} = O\left(\frac{n \lg n + n \lg \frac{1}{\varepsilon}}{\delta \cdot \lg \frac{1}{\varepsilon}}\right) \xrightarrow{\varepsilon \rightarrow 0} t_{\text{rel}} = O\left(\frac{n}{\delta}\right)$$

HW: What happens for Dobrushin++?

$$\lambda_{\text{max}}(I) \leq 1 - \delta \Rightarrow t_{\text{rel}} = O(\text{?})$$

## Intro to Functional Analysis

Question: Can we bound  $t_{\text{mix}}$  by  $t_{\text{rel}}$ ?

Idea: Contraction of proxy for  $d_{TV}$

Def: (f-Divergence) convex function  $f$ :

$$D_f(\nu \parallel \mu) := \underbrace{\mathbb{E}_{x \sim \mu} [f\left(\frac{\nu(x)}{\mu(x)}\right)]}_{\text{proxy for } d_{TV}(\nu, \mu)} - \underbrace{f\left(\mathbb{E}_{x \sim \mu} \left[\frac{\nu(x)}{\mu(x)}\right]\right)}_{1 \text{ when } \nu \text{ is dist}}$$

Note: We define it for any  $\nu, \mu: \Omega \rightarrow \mathbb{R}$  where  $\nu/\mu$  takes values in domain of  $f$ .

Fact:  $f$  convex  $\Rightarrow D_f(\nu \parallel \mu) \geq 0$

Proof: Jensen's inequality!  $\square$

Remark:  $f$  strongly convex  $\Rightarrow D_f = 0 \Leftrightarrow \nu = \text{const. } \mu$ .

Fact: Suppose  $N \in \mathbb{R}_{\geq 0}^{2 \times 2}$  is row stochastic.  
 ↓  
 noise or Markov chain

Then  $D_f(\nu N \parallel \mu N) \leq D_f(\nu \parallel \mu)$   
 Data processing ineq.

Proof: We apply Jensen's again:  $\pi(x, y) := \mu(x)N(x, y)$   
 ↓  
 $\pi(x|y)$

$$\frac{\nu N(y)}{\mu N(y)} = \frac{\sum_x \nu(x)N(x, y)}{\sum_{x'} \mu(x')N(x', y)} = \sum_x \frac{\pi(x, y)}{\sum_{x'} \pi(x', y)} \frac{\nu(x)}{\mu(x)}$$

$$\Rightarrow f\left(\frac{\nu N(y)}{\mu N(y)}\right) \leq \sum_x \pi(x|y) f\left(\frac{\nu(x)}{\mu(x)}\right) \Rightarrow$$

$$\mathbb{E}_{y \sim \mu N} \left[ f\left(\frac{\nu N(y)}{\mu N(y)}\right) \right] \leq \sum_x \underbrace{\mathbb{E}_{y \sim \mu N} [\pi(x|y)]}_{\substack{\text{x-marginal of } \pi, \\ \text{i.e. } \mu(x)}} f\left(\frac{\nu(x)}{\mu(x)}\right) =$$

$$\mathbb{E}_{x \sim \mu} \left[ f\left(\frac{\nu(x)}{\mu(x)}\right) \right]$$

Note that

$$\mathbb{E}_{y \sim \mu N} \left[ \frac{\nu N(y)}{\mu N(y)} \right] = \sum_y \nu N(y) = \sum_x \nu(x) \sum_y N(x, y) =$$

$$\mathbb{E}_{x \sim \mu} \left[ \frac{\nu(x)}{\mu(x)} \right]$$

Therefore  $D_f(\nu \parallel \mu) \geq D_f(\nu N \parallel \mu N) \quad \square$

Remark: This is extremely useful for Markov chains constructed as  $P = NN^0$ .

Contraction of either  $N$  or  $N^0 \Rightarrow$   
 contraction of  $P$

In lots of scenarios  $N$  by itself easier to analyze!

Remark: We usually want  $f$ s where  $d_{TV} \leq \text{some func of } D_f$

Popular Choice 1: ( $f(x) = x^2$ )  $\mathcal{X}^2$ -divergence

$$D_f(\nu \parallel \mu) = \mathcal{X}^2(\nu \parallel \mu)$$

Also called  $\text{Var}_\mu \left[ \frac{\nu}{\mu} \right]$   
variance

- Note that  $\nu$  can take  $< 0$  values too here.

- Alternative formula:

$$\mathcal{X}^2(\nu \parallel \mu) = \mathbb{E}_{x \sim \mu} \left[ \left( \frac{\nu(x)}{\mu(x)} - \mathbb{E}_{x \sim \mu} \left[ \frac{\nu(x')}{\mu(x')} \right] \right)^2 \right]$$

Thm: For  $\nu$  a dist, we have

$$d_{TV}(\nu, \mu) \leq \frac{1}{2} \sqrt{\mathcal{X}^2(\nu \parallel \mu)}$$

Proof: Apply Cauchy-Schwarz:

$$2d_{TV}(\nu, \mu) = \sum_x |\nu(x) - \mu(x)| = \mathbb{E}_{x \sim \mu} \left[ \left| \frac{\nu(x)}{\mu(x)} - 1 \right| \right] \leq \sqrt{\mathcal{X}^2(\nu \parallel \mu)} \quad \square$$

Popular Choice 2: ( $f(x) = x \log x$  for  $x \geq 0$ )

$$D_f(\nu \parallel \mu) = D_{KL}(\nu \parallel \mu)$$

$\rightarrow$  KL-divergence

Also called  $\text{Ent}_\mu \left[ \frac{\nu}{\mu} \right]$  entropy

- Note that  $\nu$  can take  $\geq 0$  values but doesn't need to sum to 1.

- Formula for dist  $\nu$ :  $D_{KL}(\nu \parallel \mu) = \mathbb{E}_{x \sim \nu} \left[ \log \left( \frac{\nu(x)}{\mu(x)} \right) \right]$

Thm ( Pinsker's ): For  $\nu$  a dist, we have

$$d_{TV}(\nu, \mu) \leq \sqrt{\frac{1}{2} D_{KL}(\nu \parallel \mu)}$$

Proof: Define noise operator that maps  $x$  to  $0$  if  $\nu(x) \geq \mu(x)$  and  $1$  if  $\nu(x) < \mu(x)$ .

$$d_{TV}(\nu N, \mu N) = d_{TV}(\nu, \mu), \quad D_{KL}(\nu N \parallel \mu N) \leq D_{KL}(\nu \parallel \mu)$$

$\Rightarrow$  Enough to prove on  $\Omega = \{0, 1\}$

The rest on HW.

Functional Analysis:

Show contraction of some  $D_f$ :

$$D_f(\nu P \parallel \mu) = D_f(\nu P \parallel \mu P) \leq (1-\rho) D_f(\nu \parallel \mu)$$

$$\Rightarrow D_f(\nu P^t \parallel \mu) \leq (1-\rho)^t D_f(\nu \parallel \mu)$$

Variance /  $\chi^2$

$$t \gtrsim \frac{\lg(\chi^2(\nu_0 \parallel \mu) / \varepsilon^2)}{\rho}$$

$$\Rightarrow d_{TV}(\nu_0 P^t, \mu) \leq \varepsilon$$

How large can  $\chi^2(\nu_0 \parallel \mu)$

be?  $\frac{\nu_0(x)}{\mu(x)} \leq \frac{1}{\mu_{\min}}$

$$t_{\text{mix}} = O\left(\frac{\lg\left(\frac{1}{\mu_{\min}}\right)}{\rho}\right)$$

Entropy /  $D_{KL}$

$$t \gtrsim \frac{\lg(D_{KL}(\nu_0 \parallel \mu) / \varepsilon^2)}{\rho}$$

$$\Rightarrow d_{TV}(\nu_0 P^t, \mu) \leq \varepsilon$$

How large can  $D_{KL}(\nu_0 \parallel \mu)$

be?  $\lg \frac{\nu_0(x)}{\mu(x)} \leq \lg \frac{1}{\mu_{\min}}$

$$t_{\text{mix}} = O\left(\frac{\lg \lg\left(\frac{1}{\mu_{\min}}\right)}{\rho}\right)$$

These inequalities are related to

Poincare & MLSI

future: cont. time

Spectral gap  $\Leftrightarrow$  Contraction of  $\chi^2$

Suppose  $N \in \mathbb{R}_{>0}^{\Omega \times \Omega}$  is row-stochastic and  $\mu$  dist on  $\Omega$ . Let  $\mu' = \mu N$  and  $N^0$  be time-reversal of  $N$  w.r.t  $\mu$ .

$D_\mu$ : diagonal matrix with  $\mu$  on diag.  $D_{\mu'}$ : diagonal matrix with  $\mu'$  on diag.

$$D_{\mu'} N^0 = (D_\mu N)^T = N^T D_\mu$$

Formula for  $\chi^2$ : Suppose  $\nu$  sums to 0.

$$\chi^2(\nu \parallel \mu) = \|\nu D_\mu^{-\frac{1}{2}}\|_2^2$$

For such a  $v$  we have  $(vN)^t = 0$  too,  
 so  $\chi^2(vN \| \mu N) = \|vND_{\mu'}^{-\frac{1}{2}}\|_2^2$ .

What is

$$\sup \left\{ \frac{\|vND_{\mu'}^{-\frac{1}{2}}\|_2^2}{\|vD_{\mu}^{-\frac{1}{2}}\|_2^2} \mid v1^t = 0 \right\} ?$$

This is equal to  $\lambda_2(NN^0)$ .

[if time ...]

Contraction for arbitrary  $v \leftrightarrow$  Contraction for  $v1^t = 0$

Variance =  $1 - \lambda_2(NN^0)$

Corollary: Time-reversible:

$$\lambda_2(P^2) = \max(\lambda_2(N)^2, \lambda_1(N)^2)$$

Proof: Let  $v = vD_{\mu}^{-1/2}$  ← change of variable

Let  $A = D_{\mu}^{1/2}ND_{\mu'}^{-1/2}$ . Then

$$\sup \left\{ \frac{vAA^T v^T}{v v^T} \mid vD_{\mu}^{1/2}1^t = 0 \right\}.$$

Note that

$$AA^T = D_{\mu}^{1/2}ND_{\mu'}^{-1}N^T D_{\mu}^{-1/2} = D_{\mu}^{1/2} \underbrace{NN^0}_{N^0} D_{\mu}^{-1/2}$$

is a symmetric matrix and similar to  $NN^0$ , so has same eigenvalues.

$$AA^T D_{\mu}^{1/2}1^t = D_{\mu}^{1/2}1^t \leftarrow \text{eigenvector for eigenvalue 1.}$$

So we are looking at

$$\sup \left\{ \frac{vAA^T v}{v v^T} \mid v \text{ orthogonal to top eigenvector} \right\}$$

$$= \lambda_2(AA^T) = \lambda_2(NN^0) \text{ 😊}$$

# Fourier Analysis

Finite Abelian group:  $G = \mathbb{Z}_{n_1} \times \mathbb{Z}_{n_2} \times \dots \times \mathbb{Z}_{n_k}$   
 with  $+$ : component wise addition.  $\uparrow$  mod

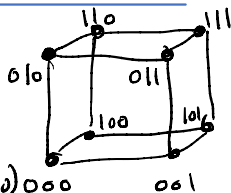
Take dist  $\pi$  over  $G$ .  
 $\uparrow$  think sparse support.

We get an Abelian walk:

$$X \mapsto X + Z \quad \leftarrow \begin{array}{l} \text{random sample} \\ \text{from } \pi \end{array}$$

Example.  $G = \mathbb{Z}_2^n$

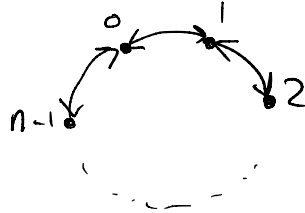
$\pi$ :  $\begin{cases} \text{w.p. } \frac{1}{2} \text{ choose } (0, \dots, \rightarrow 0) \\ \text{w.p. } \frac{1}{2n} \text{ choose } e_i = (0, \dots, \rightarrow 1, \dots, 0) \end{cases}$



This is Glauber dynamics.

Example. (cycle)

$\pi$ :  $\begin{cases} \text{w.p. } \frac{1}{2} \text{ choose } +1 \\ \text{w.p. } \frac{1}{2} \text{ choose } -1 \end{cases}$



Fact 1:  $\mu = \text{uniform}$  always stationary

Fact 2:  $\pi$  symmetric  $\iff$  time-reversible  
 $\downarrow$   
 $\pi(z) = \pi(-z)$

Fact 3: support  $\pi$  generate  $G \iff$  irreducible.

Thm: Regardless of  $\pi$ , eigenvectors are fixed; they are characters of  $G$ .

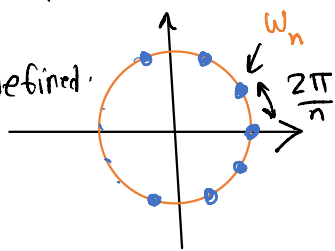
Def: (Character):  $\chi: G \rightarrow \mathbb{C}^*$  s.t.  
 $\chi(x+y) = \chi(x)\chi(y)$ .



Example:  $(\mathbb{Z}_n)$ :

Pick  $\omega \in \mathbb{C}$  s.t.  $\omega^n = 1$ .

Then  $\chi(x) = \omega^x$  is well-defined.



All characters are

$$\chi(x) = \exp\left(\frac{2\pi i \cdot tx}{n}\right) = \omega_n^{tx}$$

Characters of  $\mathbb{Z}_{n_1} \times \dots \times \mathbb{Z}_{n_k}$

$$\chi(x_1, \dots, x_k) = \omega_{n_1}^{t_1 x_1} \dots \omega_{n_k}^{t_k x_k}$$

There are exactly  $n_1 \dots n_k$  such characters.

Thm: If  $P$  is an Abelian walk, every character  $\chi$  is an eigenvector.

$$(\chi P)(y) = \sum_{x \in G} \chi(x) P(x, y) =$$

$$\sum_{x \in G} \chi(x) \pi(y-x) = \sum_z \pi(z) \chi(y-z)$$
$$= \underbrace{\left( \sum_z \pi(z) \chi(z)^{-1} \right)}_{\text{Corresponding eigenvalue}} \chi(y).$$

Since there are  $n_1 \dots n_k$  many, these are an eigenbasis!

$$\mathbb{E} \left[ \chi(-z) \right]_{z \sim \pi}$$