## Problem Set 1

CS 260: Geometry of Polynomials in Algorithm Design

https://nimaanari.com/cs260-winter2020/

## 1 Real Stability vs. Log-Concavity

Let  $g(z_1, z_2) = a + bz_1 + cz_2 + dz_1z_2$  be a multi-affine bivariate polynomial with  $a, b, c, d \in \mathbb{R}_{\geq 0}$ . Prove that

- 1) *g* is real stable if and only if  $ad \leq bc$ .
- 2) log *g* is concave over  $\mathbb{R}^2_{>0}$  if and only if  $ad \leq 2bc$ .

## 2 Approximating Higher Coefficients

Let  $g(z_1, ..., z_n)$  be a real stable polynomial with nonnegative coefficients, and let  $k_1, ..., k_n \in \mathbb{Z}_{\geq 0}$ . Let

$$\alpha := \inf_{z_1, \dots, z_n > 0} \frac{g(z_1, \dots, z_n)}{z_1^{k_1} \cdots z_n^{k_n}}$$

and let  $\beta$  be the coefficient of the  $z_1^{k_1} \cdots z_n^{k_n}$  term in g. Show  $\alpha$  approximates  $\beta$  by following these steps:

- 1) Prove that  $\alpha \geq \beta$ .
- 2) Suppose n = 1, and  $g(z) = e^{cz}$  for some  $c \in \mathbb{R}_{\geq 0}$ . Note that g is not a polynomial, but you can think of it in terms of its Taylor series

$$g(z) = 1 + cz + \frac{c^2}{2}z^2 + \dots + \frac{c^k}{k!}z^k + \dots$$

The coefficient of  $z^k$  in the above expression is  $c^k/k!$ . What is  $\inf_{z>0} g(z)/z^k$ ? Show that the ratio between these two quantities  $(\alpha/\beta)$  is only a function of *k* and not *c*. Let this quantity be f(k).

- 3) [Bonus, no need to turn in solution to this part] Show that  $f(k) = O(\sqrt{k})$ .
- 4) Now let n = 1 and g(z) a real-rooted polynomial

$$g(z) = c_0 + c_1 z + c_2 z^2 + \dots + c_d z^d.$$

We showed in class that  $\{c_i/\binom{d}{i}\}_{i=0}^d$  is a log-concave sequence. Show that this implies  $\{\partial_z^i g(0)\}_{i=0}^d$  is a log-concave sequence as well.

- 5) Let  $r, s \in \mathbb{R}_{\geq 0}$  and define  $h(z) = re^{sz}$ . What quality does the sequence  $\{\partial_z^i h(0)\}_{i=0}^{\infty}$  have? Show that for some choice of r and s we will have  $\partial_z^i h(0) \geq \partial_z^i g(0)$  for all i and that equality happens for i = k, i.e.,  $\partial_z^k h(0) = \partial_z^k g(0)$ .
- 6) Conclude that for n = 1,  $f(k) \cdot \beta \ge \alpha$ .
- 7) Prove using induction that for general *n*

$$f(k_1)f(k_2)\cdots f(k_n)\beta \geq \alpha.$$

## 3 Monomer-Dimer Correlations

Let G = (V, E) be a graph with an edge weight function  $w : E \to \mathbb{R}_{\geq 0}$  and vertex weight function  $\lambda : V \to \mathbb{R}_{\geq 0}$ . For a matching  $M \subseteq E$ , we call the vertices v not adjacent to M the monomers of M.

Let  $\mu$  be the monomer distribution defined on subsets of *V*, where

$$\mu(S) \propto \sum_{M \text{ matching with monomer set } S} \prod_{v \in S} \lambda(v) \prod_{e \in M} w(e).$$

Let *g* be the following polynomial encoding  $\mu$  in its coefficients.

$$g(z_1,\ldots,z_n)=\sum_{S}\mu(S)\prod_{i\in S}z_i.$$

We showed in class that for  $z_1, \ldots, z_n \in \mathbb{C}$  with  $\operatorname{Re}(z_1), \ldots, \operatorname{Re}(z_n) > 0$  we have  $g(z_1, \ldots, z_n) \neq 0$ .

1) Let  $T \subseteq V$  and define the random variable *X* by sampling *S* according to  $\mu$  and computing

$$X := |V| + |S \cap T| - |S \cap (V - T)|.$$

Show that *X* takes only even nonnegative integer values.

2) Let h(z) be the following polynomial

$$h(z) := \mathbb{P}[X=0]z^0 + \mathbb{P}[X=2]z^2 + \dots + \mathbb{P}[X=2k]z^{2k} + \dots$$

Express *h* in terms of *g*. Show that *h* has roots only on the imaginary axis of the complex plane.

- 3) Show that  $h(\sqrt{z})$  is a real-rooted polynomial in *z*. Conclude that X/2 is distributed as a sum of independent Bernoulli random variables.
- 4) Let  $v \in V$ , and look at the conditional distribution of X/2 conditioned on  $v \in S$ . Show that this is also a sum of independent Bernoullis. [Hint: vary the weight  $\lambda(v)$ ]
- 5) Show that any mixture of the distribution of X/2 and the conditional distribution of X/2 (conditioned on  $v \in S$ ) is also a sum of independent Bernoullis. Conclude that

$$|\mathbb{E}[X] - \mathbb{E}[X \mid v \in S]| \le 2.$$

6) Let  $C \in \mathbb{R}^{n \times n}$  be a matrix defined by

$$C_{vu} = \mathbb{P}[u \in S \mid v \in S] - \mathbb{P}[u \in S].$$

Show that the  $\ell_1$  norm of the *v*-th row of *C* is exactly the quantity  $|\mathbb{E}[X] - \mathbb{E}[X | v \in S]|$  for some choice of *T*. Conclude that the  $\ell_1$  norm of rows of *C* are bounded by 2.